

## Exercises for Mathematical Logic (7 Jan 2025)

A formula  $\varphi(x)$  represents a set  $X \subseteq \mathbb{N}$  in a theory  $T$  if  $T \vdash \varphi(\bar{n})$  for all  $n \in X$ , and  $T \vdash \neg\varphi(x)$  for all  $n \in \mathbb{N} \setminus X$ .

A formula  $\varphi(x, y)$  represents in  $T$  a partial function  $f: \mathbb{N} \rightarrow \mathbb{N}$  if  $T \vdash \forall y (\varphi(\bar{n}, y) \leftrightarrow y = \bar{m})$  for all  $n, m \in \mathbb{N}$  such that  $f(n) = m$ .

**42.** All decidable sets are  $\Sigma_1$ -representable in  $\mathbf{Q}$ .

[Hint: Starting with  $\Sigma_1$ -definitions of  $X$  and  $\mathbb{N} \setminus X$ , write a  $\Sigma_1$  formula expressing “there is a witness for  $x \in X$  smaller than any witness for  $x \notin X$ ”. Use Exer. 32 (from 10 Dec) to show that it works.]

**43.** All partial computable functions are  $\Sigma_1$ -representable in  $\mathbf{Q}$ .

[Hint: Using a  $\Sigma_1$ -definition of the graph of  $f$ , adapt the witness comparison argument from Exer. 42.]

**44.** Prove Gödel’s diagonal lemma: for every formula  $\varphi(x)$ , there exists a sentence  $\alpha$  such that  $\mathbf{Q} \vdash \alpha \leftrightarrow \varphi(\bar{\ulcorner \alpha \urcorner})$ . [Hint: Using representability of a suitable computable function, construct a formula  $\psi(x)$  such that  $\mathbf{Q} \vdash \psi(\bar{\ulcorner \chi \urcorner}) \leftrightarrow \varphi(\bar{\ulcorner \chi(\bar{\ulcorner \chi \urcorner}) \urcorner})$  for all  $\chi(x)$ .]